



Universität Freiburg  
 Institut für Informatik  
 Prof. Dr. G. Lausen  
 Michael Schmidt

Georges-Köhler Allee, Geb. 51  
 D-79110 Freiburg  
 Tel. (0761) 203-8120  
 Tel. (0761) 203-8127

**Formal Foundations of Information Systems**  
**Summerterm 2009**  
 30.06.2009

## 7. Exercise Set: Petri Nets: P/T Nets

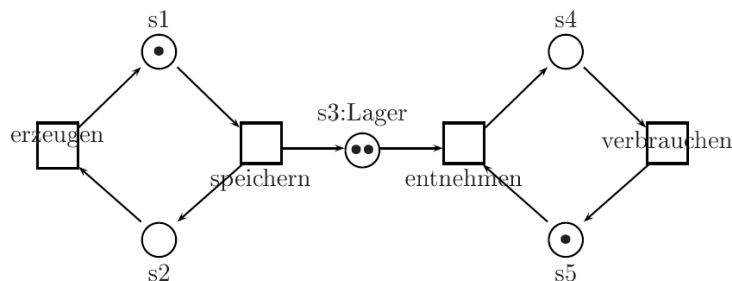
### Exercise 26 (Erreichbarkeit, Überdeckung und Beschränktheit, 1+1+1+1=4 Punkte)

Let  $N := (P, T, E, V, m_0)$  be a P/T-Net and  $m$  a marking. Prove or disprove the following two claims.

- If  $m$  is reachable in  $N$ , then  $m$  is also coverable in  $N$ .
- If  $m$  is coverable in  $N$ , then  $m$  is also reachable in  $N$ .
- If  $N$  is bounded (in german: *beschränkt*), then its reachability graph  $EG(N)$  is finite.
- If  $N$  is bounded, then the set of all words  $L_N(m_0)$  that can fire on  $N$  is finite.

### Exercise 27 (Beschränktheit, 1+2=3 Punkte)

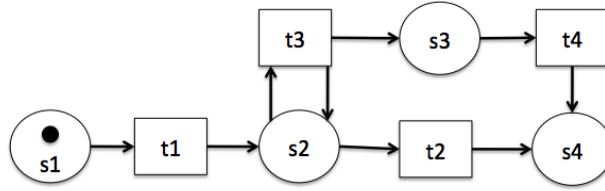
The following Petri Net  $N$  describes a producer-consumer setting. The producer creates a product that is stored in a warehouse and the consumer consumes it by taking it out of the warehouse.



- Describe this Petri Net. What are the possible markings that are reachable? How many producers and consumers are modeled in this net?
- The warehouse can store infinitely many elements, which is not applicable in a real world scenario. Therefore change  $P$  to  $P'$  (by insertion of places and transitions), so that the warehouse does not hold more than  $k$  elements at any time. Prove that  $P'$  is bounded.

### Exercise 28 (Beschränkte Stellen und Überdeckungsgraph, 5 Punkte)

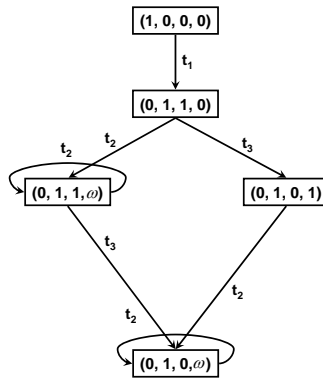
Consider the following P/T-Net:



Identify all bounded places. Use the coverage graph to solve this exercise. Also provide the help graph used to construct the coverage graph.

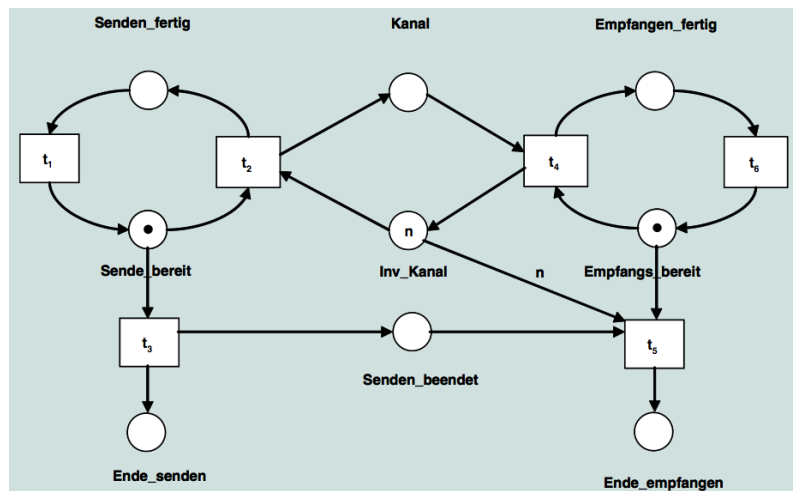
**Exercise 29 (Überdeckungsgraph, 3 Punkte)**

Model one possible Petri Net for the following coverage graph.



**Exercise 30 (Motivation Invarianten, Bonusaufgabe: 5 Punkte)**

Consider the following P/T-Net:



Try to prove or disprove the following claims using the tools from the lecture, in particular using the reachability graph. Which claims require new techniques to be solved properly?

- a) At most  $n$  messages can be transmitted in parallel.
- b) The sender is either in his final state (*Ende\_senden*), ready (*Senden\_bereit*), or finished (*Senden\_fertig*).
- c) The receiver is either in his final state (*Ende\_empfangen*), ready (*Empfangs\_bereit*), or finished (*Empfangen\_fertig*).
- d) The final state of the receiver can be reached if and only if the channel is empty and the sender has reached his final state.

Due by: 07.07.2009